

Bifibrational functorial semantics of parametric polymorphism

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Abstract

Reynolds' theory of parametric polymorphism captures the invariance of polymorphically typed programs under change of data representation. Semantically, reflexive graph categories and fibrations are both known to give a categorical understanding of parametric polymorphism. This paper contributes further to this categorical perspective by showing the relevance of bifibrations. We develop a bifibrational framework for models of System F that are *parametric*, in that they verify the Identity Extension Lemma and Reynolds' Abstraction Theorem. We also prove that our models satisfy expected properties, such as the existence of initial algebras and final coalgebras, and that parametricity implies dinaturality.

Keywords: Parametricity, logical relations, System F, fibred category theory.

1 Introduction

Strachey [30] called a polymorphic function *parametric* if its behaviour is uniform across all of its type instantiations. Reynolds [25] made this mathematically precise by formulating the notion of *relational parametricity*, in which the uniformity of parametric polymorphic functions is captured by requiring them to preserve all logical relations between instantiated types. Relational parametricity has proven to be a key technique for formally establishing properties of software systems, such as representation independence [1,6], equivalences between programs [15], and useful (“free”) theorems about programs from their types alone [31]. In this paper, we treat relational parametricity for the polymorphic λ -calculus System F [10], which forms the core of many modern programming languages and verification systems. Hermida, Reddy, and Robinson [14] give a good introduction to relational parametricity.

Since category theory underpins and informs many of the key ideas underlying modern programming languages, it is natural to ask whether it can provide a useful perspective on parametricity as well. Ma and Reynolds [19] developed the first categorical formulation of relational parametricity, but their models were complicated

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and challenging to understand. Moreover, Birkedal and Rosolini discovered that not all expected consequences of parametricity necessarily hold in their models (see [4]).

Another line of work, begun by O’Hearn and Tennent [21] and Robinson and Rosolini [28], and later refined by Dunphy and Reddy [7], uses reflexive graphs to model relations and functors between reflexive graph categories to model types. This is the state of the art for functorial semantics for parametric polymorphism. Interpreting types as functors is conceptually elegant and Dunphy and Reddy show that this framework is powerful enough to prove expected results, such as the existence of initial algebras for strictly positive type expressions [5]. However, since reflexive graph categories are relatively unknown mathematical structures, much of this development has had to be carried out from scratch. We propose to instead take the more established fibrational view of logic from the outset, and thus to analyse parametricity through the powerful lens of categorical type theory [16].

In doing so, we follow an extensive line of work by Hermida [12,13] and Birkedal and Møgelberg [4], who use fibrations to construct sophisticated categorical models not only of parametricity, but also of its logical structure in terms of Abadi-Plotkin logic [24]. Abadi-Plotkin logic is a formal logic for parametric polymorphism that includes predicate logic *and* a polymorphic lambda calculus, and thus requires significant machinery to handle. Using this machinery, Birkedal and Møgelberg are able to go beyond Dunphy and Reddy’s results and, for instance, prove that all positive type expressions — not just the strictly positive ones as for Dunphy and Reddy — have initial algebras. However, these impressive results come at the price of the complexity of the notions involved. Our aim is to achieve the same results in a simpler setting¹, closer to Dunphy and Reddy’s functorial semantics. We end up with a notion of model in which each type is interpreted as an equality preserving fibred functor and each term is interpreted as a fibred natural transformation. This is quite similar to the models produced by the parametric completion process of Robinson and Rosolini [28] (see also Birkedal and Møgelberg [4, Section 8]) and to Mitchell and Scedrov’s relator model [20], but with a more general notion of relation given by a fibration. We thus combine the generality of Birkedal and Møgelberg’s fibrational models with the simplicity of Dunphy and Reddy’s functorial semantics.

Our central innovation is the use of *bifibrations* to achieve this “sweet spot” in the study of parametricity. This is not necessary for the *definition* of our framework, for which Lawvere equality [17] (i.e., opreindexing along diagonals only) suffices, but it helps considerably with both the concrete interpretation of \forall -types [9] and the handling of graph relations. At a technical level, our strongest result is to use our simpler framework to recover all the expected consequences of parametricity that Birkedal and Møgelberg [4] prove using Abadi-Plotkin logic. In particular, we go beyond Dunphy and Reddy’s result by deriving, this time with a functorial semantics, initial algebras for all positive type expressions, rather than just for strictly positive ones. Nevertheless, this paper is in no way intended as the final word on fibrational parametricity. Instead, we hope the simple re-conceptualization of parametricity we offer here — replacing the usual categorical interpretations of types as functors and

¹ We stress again that we are not trying to model all of Abadi-Plotkin logic, but rather only type systems involving parametric polymorphism. Indeed, with respect to Abadi-Plotkin logic, we could not hope to improve upon the results of Birkedal and Møgelberg [4], who give a sound and complete semantics.

terms as natural transformations with their fibred counterparts — will open the way to the study of parametricity in richer settings, e.g., proof-relevant ones.

Structure of the paper: In Section 2 we give a short introduction to bifibrations. We recall Reynolds’ relational interpretation of System F, the Identity Extension Lemma and the Abstraction Theorem in Section 3. We then extract bifibrational generalisations of these in Section 4, and construct our parametric models. In Section 5 we show that our models behave as expected by deriving initial algebras for all definable functors and proving that parametricity implies (di)naturality. Finally, we instantiate our framework to derive both “standard” and new models of relational parametricity in Section 6. Section 7 concludes and discusses future work.

2 A Fibrational Toolbox for Relational Parametricity

We give a brief introduction to fibrations; more details can be found in, e.g., [16].

Definition 2.1 Let $U : \mathcal{E} \rightarrow \mathcal{B}$ be a functor. A morphism $g : Q \rightarrow P$ in \mathcal{E} is *cartesian* over $f : X \rightarrow Y$ in \mathcal{B} if $Ug = f$ and, for every $g' : Q' \rightarrow P$ in \mathcal{E} with $Ug' = f \circ v$ for some $v : UQ' \rightarrow X$, there exists a unique $h : Q' \rightarrow Q$ with $Uh = v$ and $g' = g \circ h$. A morphism $g : P \rightarrow Q$ in \mathcal{E} is *opcartesian* over $f : X \rightarrow Y$ in \mathcal{B} if $Ug = f$ and, for every $g' : P \rightarrow Q'$ in \mathcal{E} with $Ug' = v \circ f$ for some $v : Y \rightarrow UQ'$, there exists a unique $h : Q \rightarrow Q'$ with $Uh = v$ and $g' = h \circ g$.

We write f_P^\S for the cartesian morphism over f with codomain P , and f_\S^P for the opcartesian morphism over f with domain P . Such morphisms are unique up to isomorphism. If P is an object of \mathcal{E} then we write f^*P for the domain of f_P^\S and $\Sigma_f P$ for the codomain of f_\S^P .

Definition 2.2 A functor $U : \mathcal{E} \rightarrow \mathcal{B}$ is a *fibration* if for every object P of \mathcal{E} and every morphism $f : X \rightarrow UP$ in \mathcal{B} , there is a cartesian morphism $f_P^\S : Q \rightarrow P$ in \mathcal{E} over f . Similarly, U is an *opfibration* if for every object P of \mathcal{E} and every morphism $f : UP \rightarrow Y$ in \mathcal{B} , there is an opcartesian morphism $f_\S^P : P \rightarrow Q$ in \mathcal{E} over f . A functor U is a *bifibration* if it is both a fibration and an opfibration.

If $U : \mathcal{E} \rightarrow \mathcal{B}$ is a fibration, opfibration, or bifibration, then \mathcal{E} is its *total category* and \mathcal{B} is its *base category*. An object P in \mathcal{E} is *over* its image UP and similarly for morphisms. A morphism is *vertical* if it is over id . We write \mathcal{E}_X for the *fibred category* over an object X in \mathcal{B} , i.e., the subcategory of \mathcal{E} of objects over X and morphisms over id_X . For $f : X \rightarrow Y$ in \mathcal{B} , the function mapping each object P of \mathcal{E} to f^*P extends to a functor $f^* : \mathcal{E}_Y \rightarrow \mathcal{E}_X$ mapping each morphism $k : P \rightarrow P'$ in \mathcal{E}_Y to the morphism f^*k with $kf_P^\S = f_{P'}^\S f^*k$. The universal property of $f_{P'}^\S$ ensures the existence and uniqueness of f^*k . We call f^* the *reindexing functor along f* . A similar situation holds for opfibrations; the functor $\Sigma_f : \mathcal{E}_X \rightarrow \mathcal{E}_Y$ extending the function mapping each object P of \mathcal{E} to $\Sigma_f P$ is the *opreindexing functor along f* .

We write $|\mathcal{C}|$ for the discrete category of \mathcal{C} . If $U : \mathcal{E} \rightarrow \mathcal{B}$ is a functor, then the *discrete functor* $|U| : |\mathcal{E}| \rightarrow |\mathcal{B}|$ is induced by the restriction of U to $|\mathcal{E}|$. If $n \in \mathbb{N}$, then \mathcal{E}^n denotes the n -fold product of \mathcal{E} in Cat . The *n -fold product of U* , denoted $U^n : \mathcal{E}^n \rightarrow \mathcal{B}^n$, is the functor defined by $U^n(X_1, \dots, X_n) = (UX_1, \dots, UX_n)$.

Lemma 2.3 *If $U : \mathcal{E} \rightarrow \mathcal{B}$ is a functor then $|U| : |\mathcal{E}| \rightarrow |\mathcal{B}|$ is a bifibration. If U is a (bi)fibration then so is $U^n : \mathcal{E}^n \rightarrow \mathcal{B}^n$ for any natural number n . \square*

To formulate Reynolds’ relational parametricity categorically, we define the category \mathbf{Rel} of relations over \mathbf{Set} and the relations fibration on \mathbf{Set} [16].

Definition 2.4 The category \mathbf{Rel} has triples (A, B, R) as objects, where A, B , and R are sets and $R \subseteq A \times B$. A morphism $(A, B, R) \rightarrow (A', B', R')$ is a pair (f, g) , where $f : A \rightarrow A'$ and $g : B \rightarrow B'$, such that if $(a, b) \in R$ then $(fa, gb) \in R'$.

We write (A, B, R) as just R when A and B are immaterial or clear from context. Note that \mathbf{Rel} is *not* the category whose objects are sets and whose morphisms are relations, which also sometimes appears in the literature. Each set A has an associated *equality relation* defined by $\mathbf{Eq} A = \{(a, a) \mid a \in A\}$.

Example 2.5 The functor $U : \mathbf{Rel} \rightarrow \mathbf{Set} \times \mathbf{Set}$ sending (A, B, R) to (A, B) is called the *relations fibration on \mathbf{Set}* . To see that U is indeed a fibration, let $(f, g) : (X_1, X_2) \rightarrow (Y_1, Y_2)$ be a morphism in $\mathbf{Set} \times \mathbf{Set}$ with $UR = (Y_1, Y_2)$ for some R in \mathbf{Rel} . If we define $(f, g)^*R \subseteq X_1 \times X_2$ by $(x_1, x_2) \in (f, g)^*R$ iff $(fx_1, gx_2) \in R$, then (f, g) is a cartesian morphism from $(f, g)^*R$ to R over (f, g) . It is also easy to see that U is an opfibration, with opreindexing given by forward image. Thus, U is a bifibration. We denote the fibre over (A, B) in the relations fibration on \mathbf{Set} by $\mathbf{Rel}(A, B)$.

Definition 2.6 Let $U : \mathcal{E} \rightarrow \mathcal{B}$ and $U' : \mathcal{E}' \rightarrow \mathcal{B}'$ be fibrations. A *fibred functor* $F : U \rightarrow U'$ comprises two functors $F_0 : \mathcal{B} \rightarrow \mathcal{B}'$ and $F_1 : \mathcal{E} \rightarrow \mathcal{E}'$ such that $U'F_1 = F_0U$ and cartesian morphisms are preserved, i.e., if f is cartesian in \mathcal{E} over g in \mathcal{B} then F_1f is cartesian in \mathcal{E}' over F_0g in \mathcal{B}' . If $F' : U \rightarrow U'$ is another fibred functor, then a *fibred natural transformation* $\eta : F \rightarrow F'$ comprises two natural transformations $\eta_0 : F_0 \rightarrow F'_0$ and $\eta_1 : F_1 \rightarrow F'_1$ such that $U'\eta_1 = \eta_0U$.

In this paper we use fibred functors and fibred transformations to interpret System F types and terms, and show that under mild conditions this gives parametric models.

3 Reynolds’ Model of Relational Parametricity

We now describe Reynolds’ set-theoretic model of relational parametricity: first concretely, and then in terms of the relations fibration $\mathbf{Rel} \rightarrow \mathbf{Set} \times \mathbf{Set}$. As Reynolds discovered, there are in fact no set-theoretic models if the meta-theory is classical logic [26], but the following makes sense in the (intuitionistic) internal language of a topos [22], or in the Calculus of Constructions with impredicative \mathbf{Set} . Throughout, we assume a standard syntax for System F.

3.1 Semantics of Types

Reynolds presents two “parallel” semantics for System F: a standard set-based semantics $\llbracket - \rrbracket_o$, and a relational semantics $\llbracket - \rrbracket_r$. Given $\Gamma \vdash T$ type, where the type context Γ contains $|\Gamma| = n$ type variables, Reynolds defines interpretations $\llbracket T \rrbracket_o : |\mathbf{Set}|^n \rightarrow \mathbf{Set}$ and $\llbracket T \rrbracket_r : |\mathbf{Rel}|^n(A, B) \rightarrow \mathbf{Rel}(\llbracket T \rrbracket_o A, \llbracket T \rrbracket_o B)$ by structural induction on type judgements as follows:

- Type variables: $\llbracket X_i \rrbracket_o A = A_i$ and $\llbracket X_i \rrbracket_r R = R_i$
- Arrow types:

$$\begin{aligned} \llbracket T_1 \rightarrow T_2 \rrbracket_o A &= \llbracket T_1 \rrbracket_o A \rightarrow \llbracket T_2 \rrbracket_o A \\ \llbracket T_1 \rightarrow T_2 \rrbracket_r R &= \{(f, g) \mid (a, b) \in \llbracket T_1 \rrbracket_r R \Rightarrow (fa, gb) \in \llbracket T_2 \rrbracket_r R\} \end{aligned}$$

- Forall types:

$$\begin{aligned} \llbracket \forall X.T \rrbracket_o A &= \{f : \prod_{S:\mathbf{Set}} \llbracket T \rrbracket_o(A, S) \mid \forall R' \in \mathbf{Rel}(A', B'). (fA', fB') \in \llbracket T \rrbracket_r(\mathbf{Eq} A, R')\} \\ \llbracket \forall X.T \rrbracket_r R &= \{(f, g) \mid \forall R' \in \mathbf{Rel}(A', B'). (fA', gB') \in \llbracket T \rrbracket_r(R, R')\} \end{aligned}$$

The definitions of $\llbracket \forall X.T \rrbracket_o$ and $\llbracket \forall X.T \rrbracket_r$ depend crucially on one another. Thus, we do not really have two semantics — one based on \mathbf{Set} and one based on \mathbf{Rel} — but rather *a single semantics* based on the relations fibration $U : \mathbf{Rel} \rightarrow \mathbf{Set} \times \mathbf{Set}$. In other words, Reynolds’ definitions of $\llbracket - \rrbracket_o$ and $\llbracket - \rrbracket_r$ entail the following theorem:

Theorem 3.1 (Fibrational Semantics of Types) *Let U be the relations fibration on \mathbf{Set} . Every judgement $\Gamma \vdash T$ induces a fibred functor $\llbracket T \rrbracket : |U|^{\Gamma} \rightarrow U$.*

$$\begin{array}{ccc} |\mathbf{Rel}|^{\Gamma} & \xrightarrow{\llbracket T \rrbracket_r} & \mathbf{Rel} \\ \downarrow |U|^{\Gamma} & & \downarrow U \\ |\mathbf{Set}|^{\Gamma} \times |\mathbf{Set}|^{\Gamma} & \xrightarrow{\llbracket T \rrbracket_o \times \llbracket T \rrbracket_o} & \mathbf{Set} \times \mathbf{Set} \end{array}$$

□

Since the domain of $\llbracket T \rrbracket_r$ is a discrete category, requiring that $\llbracket T \rrbracket$ is a fibred functor amounts simply to requiring that the above diagram commutes. In particular, no preservation of cartesian morphisms by $\llbracket T \rrbracket_r$ is needed. Reynolds does not give a functorial action of types on morphisms. This is reflected in the appearance of discrete categories in Theorem 3.1. As a result, Reynolds’ pointwise interpretation of function spaces is the exponential in the functor category $|U|^{\Gamma} \rightarrow U$ [27]. How parametricity treats the action on morphisms will become clear in Section 5.1; instead of acting on morphisms, the interpretation of types act on *graph relations* induced by morphisms. For now, we simply note that the use of discrete domains does not take us out of the fibrational framework; Lemma 2.3 ensures that $\llbracket T \rrbracket$ is a functor between fibrations. The Identity Extension Lemma (IEL) is key for many applications of parametricity. It says that every relational interpretation preserves equality relations²:

Lemma 3.2 (IEL) *If $\Gamma \vdash T$ then $\llbracket T \rrbracket_r \circ |\mathbf{Eq}|^{\Gamma} = \mathbf{Eq} \circ \llbracket T \rrbracket_o$.*

□

² Reynolds’ approach also handles “identity relations” that are not equality relations, such as the information order on domains. In this paper, like many others [2, 4, 13, 24], we only treat equality relations. In future work, we hope to give an axiomatic account of “identity relations” similar to that of Dunphy and Reddy [7].

3.2 Semantics of Terms

Reynolds' main result is his Abstraction Theorem, stating that all terms send related environments to related values. Reynolds first gives set-valued and relational interpretations of term contexts $\Delta = x_1 : T_1, \dots, x_n : T_n$ by defining $\llbracket \Delta \rrbracket_o = \llbracket T_1 \rrbracket_o \times \dots \times \llbracket T_n \rrbracket_o$ and $\llbracket \Delta \rrbracket_r = \llbracket T_1 \rrbracket_r \times \dots \times \llbracket T_n \rrbracket_r$. This defines a fibred functor $\llbracket \Delta \rrbracket : |U|^{|\Gamma|} \rightarrow U$. Reynolds' then interprets each judgement $\Gamma; \Delta \vdash t : T$ as a family of functions $\llbracket t \rrbracket_o : \llbracket \Delta \rrbracket_o S \rightarrow \llbracket T \rrbracket_o S$ for each environment $S \in |\text{Set}|^{|\Gamma|}$. We omit the standard definition of $\llbracket t \rrbracket_o$ here. Finally, Reynolds proves:

Theorem 3.3 (Abstraction Theorem) *Let $A, B \in \text{Set}^{|\Gamma|}$, $R \in \text{Rel}^{|\Gamma|}(A, B)$, $a \in \llbracket \Delta \rrbracket_o A$, and $b \in \llbracket \Delta \rrbracket_o B$. For every term $\Gamma; \Delta \vdash t : T$, if $(a, b) \in \llbracket \Delta \rrbracket_r R$, then $(\llbracket t \rrbracket_o A a, \llbracket t \rrbracket_o B b) \in \llbracket T \rrbracket_r R$. Or, more concisely, fibrationally: every judgement $\Gamma; \Delta \vdash t : T$ is interpreted as a fibred natural transformation $(\llbracket t \rrbracket_o, \llbracket t \rrbracket_r) : \llbracket \Delta \rrbracket \rightarrow \llbracket T \rrbracket$.*

$$\begin{array}{ccc}
 |\text{Rel}|^{|\Gamma|} & \begin{array}{c} \xrightarrow{\llbracket \Delta \rrbracket_r} \\ \Downarrow \llbracket t \rrbracket_r \\ \xrightarrow{\llbracket T \rrbracket_r} \end{array} & \text{Rel} \\
 \downarrow |U|^{|\Gamma|} & & \downarrow U \\
 |\text{Set}|^{|\Gamma|} \times |\text{Set}|^{|\Gamma|} & \begin{array}{c} \xrightarrow{\llbracket \Delta \rrbracket_o \times \llbracket \Delta \rrbracket_o} \\ \Downarrow \llbracket t \rrbracket_o \times \llbracket t \rrbracket_o \\ \xrightarrow{\llbracket T \rrbracket_o \times \llbracket T \rrbracket_o} \end{array} & \text{Set} \times \text{Set}
 \end{array}$$

□

It is worthwhile to unpack the fibrational statement of the theorem: Since the domains of the functors $\llbracket \Delta \rrbracket_o$ and $\llbracket T \rrbracket_o$ are discrete, the interpretation $\llbracket t \rrbracket_o$ actually defines a (vacuously natural) transformation $\llbracket t \rrbracket_o : \llbracket \Delta \rrbracket_o \rightarrow \llbracket T \rrbracket_o$. By the definition of morphisms in the category Rel , the existence of the (again, vacuously natural) transformation $\llbracket t \rrbracket_r$ over $\llbracket t \rrbracket_o \times \llbracket t \rrbracket_o$ is exactly the statement that if $(a, b) \in \llbracket \Delta \rrbracket_r R$, then $(\llbracket t \rrbracket_o A a, \llbracket t \rrbracket_o B b) \in \llbracket T \rrbracket_r R$ — the verbose conclusion of the theorem.

Reynolds' original formulation of the Abstraction Theorem makes it seem at first glance as though it asserts a property of $\llbracket t \rrbracket_o$. Surprisingly, however, our fibrational version makes it clear that the Abstraction Theorem actually states the existence of additional algebraic structure given by $\llbracket t \rrbracket_r$ and, more generally, the interpretation of terms as fibred natural transformations. Taking this point of view and exposing this heretofore hidden structure opens the way to our bifibrational generalisation of Reynolds' model.

4 Bifibrational Relational Parametricity

Thus far we have only shown how to view Reynolds' notion of parametricity in terms of the specific fibration $U : \text{Rel} \rightarrow \text{Set} \times \text{Set}$. We now generalise this to other fibrations. This requires that we generalise $\llbracket - \rrbracket_o$ and $\llbracket - \rrbracket_r$ in such a way that the IEL and the Abstraction Theorem hold, which in turn requires that we define equality functors for these other fibrations. The construction of equality functors is standard in any fibration with the necessary infrastructure [16], but we briefly describe it here for completeness. The first step is to note that the relations fibration from Example 2.5 arises from the subobject fibration over Set by so-called *change of base* (or pullback), and to generalise that construction.

Definition 4.1 Let $U : \mathcal{E} \rightarrow \mathcal{B}$ be a fibration and suppose \mathcal{B} has products. The fibration $\text{Rel}(U) : \text{Rel}(\mathcal{E}) \rightarrow \mathcal{B} \times \mathcal{B}$ is defined by the following change of base:

$$\begin{array}{ccc} \text{Rel}(\mathcal{E}) & \xrightarrow{q} & \mathcal{E} \\ \text{Rel}(U) \downarrow & \lrcorner & \downarrow U \\ \mathcal{B} \times \mathcal{B} & \xrightarrow{-\times-} & \mathcal{B} \end{array}$$

We call $\text{Rel}(U)$ the *relations fibration for U* , and call the objects of $\text{Rel}(\mathcal{E})$ *relations* on \mathcal{B} , to emphasise that this construction generalises the relations fibration on Set . We say that a fibration $U : \mathcal{E} \rightarrow \mathcal{B}$ has *fibred terminal objects* if each fibre \mathcal{E}_X of \mathcal{E} has a terminal object, and if reindexing preserves these terminal objects. The map sending each object X of \mathcal{B} to the terminal object in \mathcal{E}_X extends to a functor $K : \mathcal{B} \rightarrow \mathcal{E}$ called the *truth functor* for U . We can construct an equality functor for $\text{Rel}(U)$ from the truth functor for U as follows:

Definition 4.2 Let $U : \mathcal{E} \rightarrow \mathcal{B}$ be a bifibration with fibred terminal objects. If \mathcal{B} has products, then the map $X \mapsto \Sigma_{\delta_X} KX$, where δ_X is the diagonal morphism $\delta_X : X \rightarrow X \times X$, extends to the *equality functor* $\text{Eq} : \mathcal{B} \rightarrow \text{Rel}(\mathcal{E})$ for $\text{Rel}(U)$.

For this definition, it is enough to ask for opreindexing along diagonals δ_X only (this is what Birkedal and Møgelberg [4] do to model equality). When dealing with graph relations in Section 5.1, though, we will use all of the opfibrational structure to opreindex along arbitrary morphisms. Our definition specialises to the equality relation $\text{Eq } A$ when instantiated to the relations fibration on Set . The equality functor is faithful, but not always full; a counterexample is the equality functor for the identity bifibration $\text{Id} : \text{Set} \rightarrow \text{Set}$, which gives a model with *ad hoc*, rather than parametric, polymorphic functions. We thus assume in the rest of this paper that equality functors are full. This is reminiscent of Birkedal and Møgelberg’s [4] assumption that the fibration has *very strong equality*, i.e., that internal equality implies external equality, in the following sense: fullness says that if $(f, g, \alpha) : \mathbf{1} \rightarrow \text{Eq } Y$ (i.e., α shows that $f = g$ internally), then, since $\mathbf{1} = \text{Eq } \mathbf{1}_B$, $(f, g, \alpha) = (h, h, \text{Eq } h)$ for some $h : \mathbf{1}_B \rightarrow Y$ (i.e., $f = g$ externally). We use fullness of Eq at several places in Section 5 below.

We now show how to interpret arrow types and forall types as fibred functors with discrete domains. We then show that a particular class of such functors forms a $\lambda 2$ -fibration and thus a model of System F which is, in fact, parametric.

4.1 Interpreting Arrow Types

The definition of $\llbracket T \rightarrow U \rrbracket_o$ and $\llbracket T \rightarrow U \rrbracket_r$ in Section 3.1 is derived from the cartesian closed structure of Set and Rel , respectively. Moreover, the fibration $U : \text{Rel} \rightarrow \text{Set} \times \text{Set}$ preserves the cartesian closed structure, so that $\llbracket t \rrbracket_r$ is over $\llbracket t \rrbracket_o \times \llbracket t \rrbracket_o$ as required by the Abstraction Theorem. Generalising from this fibration, we can model arrow types “parametrically” — i.e., in a way satisfying the Abstraction Theorem — in any fibration $U : \mathcal{E} \rightarrow \mathcal{B}$ in which \mathcal{E} and \mathcal{B} are cartesian closed categories (CCCs) and U preserves cartesian closedness.

Definition 4.3 A fibration $U : \mathcal{E} \rightarrow \mathcal{B}$ is an *arrow fibration* if both \mathcal{E} and \mathcal{B} are CCCs, and U preserves the cartesian closed structure. A relations fibration $\text{Rel}(U)$ is an *equality preserving arrow fibration* if it is an arrow fibration and $\text{Eq} : \mathcal{B} \rightarrow \text{Rel}(\mathcal{E})$ preserves exponentials.

One advantage of working with well-studied mathematical structures such as fibrations is that many of their properties can be found in the literature. This helps in determining when a relations fibration is an equality preserving arrow fibration:

Lemma 4.4 *Let $U : \mathcal{E} \rightarrow \mathcal{B}$ be a bifibration with fibred terminal objects and \mathcal{B} be a CCC.*

- (i) *If $\text{Eq} : \mathcal{B} \rightarrow \text{Rel}(\mathcal{E})$ has a left adjoint Q , then Eq preserves exponentials iff Q satisfies the Frobenius property. Such a Q exists if $U : \mathcal{E} \rightarrow \mathcal{B}$ has full comprehension, $\text{Eq} : \mathcal{B} \rightarrow \text{Rel}(\mathcal{E})$ is full and \mathcal{B} has pushouts.*
- (ii) *If $U : \mathcal{E} \rightarrow \mathcal{B}$ is a fibred CCC and has simple products (i.e., if, for every projection $\pi_B : A \times B \rightarrow A$ in \mathcal{B} , the reindexing functor π_B^* has a right adjoint and the Beck-Chevalley condition holds), then \mathcal{E} is a CCC and U preserves the cartesian closed structure. \square*

Change of base preserves simple products and fibred structure, so $\text{Rel}(U)$ is a fibred CCC with simple products if U is. Moreover, $\mathcal{B} \times \mathcal{B}$ is a CCC if \mathcal{B} is. Lemma 4.4 thus derives structure in $\text{Rel}(U)$ from structure in U .

4.2 Interpreting Forall Types

We must generalise Reynolds' definitions of $\llbracket - \rrbracket_o$ and $\llbracket - \rrbracket_r$ for forall types to relations fibrations in such a way that the Abstraction Theorem and IEL hold. The rules for type abstraction and type application suggest that we should interpret \forall as right adjoint to weakening by a type variable. We may first try to look for such an adjoint on the base category, then another on the total category, and then try to link these adjoints. But this is the wrong idea; for the relations fibration of Example 2.5, this gives *all* polymorphic functions, not just the parametrically polymorphic ones. Instead, we require an adjoint for the combined fibred semantics. Let $|\text{Rel}(U)|^n \rightarrow_{\text{Eq}} \text{Rel}(U)$ be the category whose objects are equality preserving fibred functors from $|\text{Rel}(U)|^n$ to $\text{Rel}(U)$ and whose morphisms are fibred natural transformations between them. Then:

Definition 4.5 $\text{Rel}(U)$ is a \forall -fibration if, for every projection $\pi_n : |\text{Rel}(U)|^{n+1} \rightarrow |\text{Rel}(U)|^n$, the functor $_ \circ \pi_n : (|\text{Rel}(U)|^n \rightarrow_{\text{Eq}} \text{Rel}(U)) \rightarrow (|\text{Rel}(U)|^{n+1} \rightarrow_{\text{Eq}} \text{Rel}(U))$ has a right adjoint \forall_n and this family of adjunctions is natural in n .

We write \forall for \forall_n when n can be inferred. This definition follows, e.g., Dunphy and Reddy [7] by “baking the Identity Extension Lemma into” the definition of forall types — in the sense that the very existence of \forall requires that if F is equality preserving then so is $\forall F$ — rather than relegating it to a result to be proved *post facto*. If U is faithful, then Definition 4.5 can be reformulated in terms of more basic concepts using its opfibrational structure. The IEL then becomes a consequence of the definition, rather than an intrinsic part of it [9]. For the purposes of this paper, this abstract specification is enough.

4.3 Fibred functors with discrete domains form a parametric model

A $\lambda 2$ -fibration, i.e., a fibration $p : \mathcal{G} \rightarrow \mathcal{S}$ with fibred finite products, finite products in \mathcal{S} , fibred exponents, a generic object Ω , and simple Ω -products, is a categorical model of System F. Seely [29] gives a sound interpretation of the calculus in such fibrations. We conclude this section with the following theorem:

Theorem 4.6 *If $\text{Rel}(U)$ is an equality preserving arrow fibration and a \forall -fibration, then there is a $\lambda 2$ -fibration in which types $\Gamma \vdash T$ are interpreted as equality preserving fibred functors $\llbracket T \rrbracket : |\text{Rel}(U)|^{|\Gamma|} \rightarrow_{\text{Eq}} \text{Rel}(U)$ and terms $\Gamma; \Delta \vdash t : T$ are interpreted as fibred natural transformations $\llbracket t \rrbracket : \llbracket \Delta \rrbracket \rightarrow \llbracket T \rrbracket$. \square*

Note that Lemma 4.4 gives conditions for $\text{Rel}(U)$ to be an arrow fibration, and our other paper [9] similarly gives conditions for $\text{Rel}(U)$ to be a \forall -fibration. Unwinding the interpretation of System F in $\lambda 2$ -fibrations [29], we see that we get the following for every fibration $U : \mathcal{E} \rightarrow \mathcal{B}$ satisfying the hypotheses of the theorem: for every System F type $\Gamma \vdash T$ and term $\Gamma; \Delta \vdash t : T$, we get

- (i) a standard interpretation of $\Gamma \vdash T$ as a functor $\llbracket T \rrbracket_o : |\mathcal{B}|^{|\Gamma|} \rightarrow \mathcal{B}$;
- (ii) a relational interpretation of $\Gamma \vdash T$ as a functor $\llbracket T \rrbracket_r : |\text{Rel}(\mathcal{E})|^{|\Gamma|} \rightarrow \text{Rel}(\mathcal{E})$;
- (iii) a proof of the Identity Extension Lemma in the form of Lemma 3.2, i.e., a proof that $\llbracket T \rrbracket$ is equality preserving;
- (iv) a standard interpretation of $\Gamma; \Delta \vdash t : T$ as a natural transformation $\llbracket t \rrbracket_o : \llbracket \Delta \rrbracket_o \rightarrow \llbracket T \rrbracket_o$; and
- (v) a proof of the Abstraction Theorem in the form of Theorem 3.3, i.e., a proof that $\Gamma; \Delta \vdash t : T$ has a relational interpretation as a natural transformation $\llbracket t \rrbracket_r : \llbracket \Delta \rrbracket_r \rightarrow \llbracket T \rrbracket_r$ over $\llbracket t \rrbracket_o \times \llbracket t \rrbracket_o$.

Theorem 4.6 also gives a powerful internal language [16], where base types in type context Γ are given by fibred functors $|\text{Rel}(U)|^{|\Gamma|} \rightarrow_{\text{Eq}} \text{Rel}(U)$, and base term constants in term context Δ are given by fibred natural transformations $\llbracket \Delta \rrbracket \rightarrow \llbracket T \rrbracket$. Thus, we can use this language to reason about our models using System F. This will be used in the proofs of Theorems 5.7 and 5.11 below.

5 Consequences of parametricity

We use our new framework to derive expected consequences of parametricity. This serves as a “sanity check” for our new bifibrational conceptualisation, and shows that our framework is powerful enough to derive the same results as, e.g., Birkedal and Møgelberg [4]. At a high-level, our proof strategies are often similar to the ones found in the literature, while the proofs of individual facts are necessarily specific to our setting, and often fibrational in nature.

5.1 Graph Relations

In the fibration $U : \text{Rel} \rightarrow \text{Set} \times \text{Set}$ every function $f : X \rightarrow Y$ defines a graph relation $\langle f \rangle = \{(x, y) \mid fx = y\} \subseteq X \times Y$. This generalises to the fibrational setting, where the graph of $f : A \rightarrow B$ is obtained by reindexing the equality relation on B .

Definition 5.1 Let $U : \mathcal{E} \rightarrow \mathcal{B}$ be a fibration with fibred terminal objects and products in \mathcal{B} . The *graph* of $h : X \rightarrow Y$ in \mathcal{B} is $\langle h \rangle = (h, \text{id}_Y)^*(\text{Eq } Y)$ in $\text{Rel}(\mathcal{E})$.

The definition of $\langle h \rangle$ agrees with the set-theoretic one for the relations fibration on Set . Since reindexing preserves identities, $\langle \text{id}_A \rangle = (\text{id}_A, \text{id}_A)^*(\text{Eq } A) = \text{Eq } A$ for any object A of \mathcal{B} . In a bifibration, we can also define the graph of $f : A \rightarrow B$ in another, isomorphic way by using opfibrational structure to opreindex equality on A :

Lemma 5.2 (Lawvere [17]) *If $U : \mathcal{E} \rightarrow \mathcal{B}$ is a bifibration with fibred terminal objects that satisfies the Beck-Chevalley condition [16, Section 1.8.11], and if \mathcal{B} has products, then the graph of $h : X \rightarrow Y$ can also be described by $\langle h \rangle = \Sigma_{(\text{id}_X, h)}(\text{Eq } X)$. \square*

Being able to describe graph relations in terms of *either* reindexing *or* opreindexing in any bifibration lets us use the universal properties of both when proving theorems about them. Graph relations are the key structures that turn morphisms in \mathcal{B} into objects in $\text{Rel}(\mathcal{E})$ and, more generally, mediate the standard and relational semantics.

The *graph functor* for $\text{Rel}(U) : \text{Rel}(\mathcal{E}) \rightarrow \mathcal{B} \times \mathcal{B}$ is the functor $\langle _ \rangle : \mathcal{B}^{\rightarrow} \rightarrow \text{Rel}(\mathcal{E})$ mapping $f : X \rightarrow Y$ in \mathcal{B} to $\langle f \rangle$ in $\text{Rel}(\mathcal{E})$. To see how $\langle _ \rangle$ acts on morphisms, recall that if $f : X \rightarrow Y$ and $f' : X' \rightarrow Y'$ are objects of $\mathcal{B}^{\rightarrow}$, then a morphism from f to f' is a pair of morphisms $g : X \rightarrow X'$ and $h : Y \rightarrow Y'$ such that $f' \circ g = h \circ f$. The universal property of reindexing in $\text{Rel}(U)$ guarantees the existence of a unique morphism $\langle g, h \rangle : \langle f \rangle \rightarrow \langle f' \rangle$ over (g, h) such that the following diagram commutes:

$$\begin{array}{ccc} \langle f \rangle & \xrightarrow{(f, \text{id})^{\S}} & \text{Eq } Y \\ \exists! \langle g, h \rangle \downarrow & & \downarrow \text{Eq } h \\ \langle f' \rangle & \xrightarrow{(f', \text{id})^{\S}} & \text{Eq } Y' \end{array}$$

Lemma 5.3 *If the underlying bifibration satisfies the Beck-Chevalley condition, then $\langle _ \rangle : \mathcal{B}^{\rightarrow} \rightarrow \text{Rel}(\mathcal{E})$ is full and faithful if $\text{Eq} : \mathcal{B} \rightarrow \text{Rel}(\mathcal{E})$ is. \square*

The proof uses the opfibrational characterisation of the graph functor from Lemma 5.2. The main tool for deriving consequences of parametricity is the Graph Lemma, which relates the graph of the action of a functor on a morphism with its relational action on the graph of the morphism. Interestingly, although our setting is possibly proof-relevant (i.e., there can be multiple proofs that two elements are related), the following “logical equivalence” version of the Graph Lemma is strong enough for our applications. If $U : \mathcal{E} \rightarrow \mathcal{B}$ and $U' : \mathcal{E}' \rightarrow \mathcal{B}'$ are fibrations, we write $(F_o, F_r) : \text{Rel}(U) \rightarrow_{\text{Eq}} \text{Rel}(U')$ to indicate that functors (not necessarily fibred) $F_o : \mathcal{B} \rightarrow \mathcal{B}'$ and $F_r : \text{Rel}(\mathcal{E}) \rightarrow \text{Rel}(\mathcal{E}')$ are such that $\text{Rel}(U') \circ F_r = (F_o \times F_o) \circ \text{Rel}(U)$, and (F_o, F_r) is equality preserving, i.e., $F_r \circ \text{Eq} = \text{Eq} \circ F_o$.

Theorem 5.4 (Graph Lemma) *Assume the underlying bifibration satisfies the Beck-Chevalley condition, and let $(F_o, F_r) : \text{Rel}(U) \rightarrow_{\text{Eq}} \text{Rel}(U)$. For any $h : X \rightarrow Y$ in \mathcal{B} , there are vertical morphisms $\phi_h : \langle F_o h \rangle \rightarrow F_r \langle h \rangle$ and $\psi_h : F_r \langle h \rangle \rightarrow \langle F_o h \rangle$ in $\text{Rel}(\mathcal{E})$. \square*

Our proof of the Graph Lemma is completely independent of the specific functor (F_o, F_r) , and so in particular does not proceed by induction on the structure of

types. This is a key reason why we can go beyond Dunphy and Reddy [7] and prove the existence of initial algebras of positive, rather than just strictly positive, type expressions.

5.2 Existence of Initial Algebras

Let $F : \mathcal{C} \rightarrow \mathcal{C}$ be an endofunctor. An F -algebra is a pair (A, k_A) with A an object of \mathcal{C} and $k_A : FA \rightarrow A$ a morphism. We call A the *carrier* of the F -algebra and k_A its *structure map*. A morphism $h : A \rightarrow B$ in \mathcal{C} is an F -algebra homomorphism $h : (A, k_A) \rightarrow (B, k_B)$ if $k_B \circ (Fh) = h \circ k_A$. An F -algebra (Z, in) is *weakly initial* if, for any F -algebra (A, k_A) , there exists a mediating F -algebra homomorphism $fold[A, k_A] : (Z, in) \rightarrow (A, k_A)$. It is an *initial F -algebra* if $fold[A, k_A]$ is unique.

The literature contains other proofs that initial algebras exist in parametric models (e.g., [4,24]). Closest to our setting is Dunphy and Reddy [7], who show that strictly positive types have initial algebras. Under assumptions no stronger than theirs, we sharpen this result to *all* positive types, or, more generally, all functors on our parametric models that are strong (see below) and equality preserving.

Let $F = (F_o, F_r) : \text{Rel}(U) \rightarrow_{\text{Eq}} \text{Rel}(U)$ be a functor (note that the domain of F is not discrete and that F need not preserve cartesian morphisms) with a *strength* $t = (t_o, t_r)$, i.e., a family of morphisms $(t_o)_{A,B} : A \Rightarrow B \rightarrow F_o A \Rightarrow F_o B$ and $(t_r)_{R,S} : R \Rightarrow S \rightarrow F_r R \Rightarrow F_r S$ with $(t_r)_{R,S}$ over $((t_o)_{A,B}, (t_o)_{C,D})$ if R is over (A, B) and S is over (C, D) , such that t preserves identity and composition. A functor with a strength is said to be *strong*. Because of the discrete domains, t is a natural transformation from $_ \Rightarrow _$ to $F_ \Rightarrow F_$ in $|\text{Rel}(U)|^2 \rightarrow_{\text{Eq}} \text{Rel}(U)$, and thus $\alpha, \beta; \cdot \vdash \underline{t} : (\alpha \rightarrow \beta) \rightarrow (\underline{F}[\alpha] \rightarrow \underline{F}[\beta])$ represents the action of F on morphisms in the internal language. All type expressions with one free type variable occurring only positively give rise to strong functors, but there are further examples of such functors, for instance if the model contains non-System F type constructions with natural functorial (and relational) interpretations — for example, those of dependent types in **Set**. We will show that an initial F_o -algebra exists. For this, we first construct a weak initial F_o -algebra, which can be done in any $\lambda 2$ -fibration. Using the internal language, we define Z by $(Z_o, Z_r) = \llbracket \forall X. (\underline{F}X \rightarrow X) \rightarrow X \rrbracket$.

Lemma 5.5 Z_o is the carrier of a weak initial F_o -algebra (Z_o, in_o) with mediating morphism $fold_o[A, k]$ and Z_r is the carrier of a weak initial F_r -algebra (Z_r, in_r) with mediating morphism $fold_r[A, k]$. \square

To show that $fold_o$ is unique, we use the graph relations from Section 5.1. Recall that a category with a terminal object $\mathbf{1}$ is *well-pointed* if, for any $f, g : A \rightarrow B$, we have $f = g$ iff $f \circ e = g \circ e$ for all $e : \mathbf{1} \rightarrow A$. Like Dunphy and Reddy [7], we only consider well-pointed base categories; well-pointedness is used to convert internal language reasoning in non-empty contexts to closed contexts, so that we can apply semantic techniques such as Theorem 5.4.

Lemma 5.6 Assume that the underlying bifibration satisfies the Beck-Chevalley condition, and that **Eq** is full.

- (i) If \mathcal{B} is well-pointed, then $fold_o[Z_o, in_o] = \text{id}_Z$.

- (ii) For every F_o -algebra homomorphism $h : (Z_o, in_o) \rightarrow (A, k_A)$, we have that $h \circ fold_o[Z_o, in_o] = fold_o[A, k_A]$. \square

The proofs of the two parts of Lemma 5.6 are similar: both use the graph functor to map commuting diagrams in \mathcal{B} to morphisms in $\text{Rel}(\mathcal{E})$, and then use the Graph Lemma to see that these morphisms are F_r -algebras. Lemma 5.5 and Lemma 5.6 together now immediately imply the main result.

Theorem 5.7 *If the underlying bifibration satisfies the Beck-Chevalley condition, Eq is full, and \mathcal{B} is well-pointed, then (Z_o, in_o) is an initial F_o -algebra.* \square

We show in Section 6 that these hypothesis cannot be weakened. One may wonder if the above result can be strengthened to get not only an initial F_o -algebra, but also an initial F_r -algebra. Certainly this is possible for the relations fibration $\text{Rel} \rightarrow \text{Set} \times \text{Set}$, since relations in Rel are proof irrelevant: maps either preserve relatedness or not. This translates in the axiomatic bifibrational setting to requiring the fibration $\text{Rel}(\mathcal{E}) \rightarrow \mathcal{B} \times \mathcal{B}$ to be faithful. When it is, the weak initial F_r -algebra is, in fact, initial: faithfulness implies the required uniqueness.

5.3 Existence of final coalgebras

We can also dualise the proof from the previous section to show the existence of final coalgebras in the usual manner [11]. As usual, this requires us to first encode products and existential types in System F. We encode products as $A \times B = \forall Y.(A \rightarrow B \rightarrow Y) \rightarrow Y$. This supports the usual pairing and projection operations, as well as surjective pairing using parametricity. We encode existential types by $\exists X.T = \forall Y.(\forall X.(T \rightarrow Y)) \rightarrow Y$. We can support introduction and elimination rules

$$\frac{\Gamma \vdash A \text{ type} \quad \Gamma; \Delta \vdash u : T[A/X]}{\Gamma; \Delta \vdash \langle A, u \rangle : \exists X.T(X)} \quad \frac{\Gamma; \Delta \vdash t : \exists X.T \quad \Gamma, Z; \Delta, y : T[Z/X] \vdash s : S}{\Gamma; \Delta \vdash (\text{open } t \text{ as } \langle Z, y \rangle \text{ in } s) : S}$$

with the conversion $\text{open } \langle A, t \rangle \text{ as } \langle Z, y \rangle \text{ in } s = s[X/A, y/t]$ by defining $\langle A, t \rangle = \Lambda Y.\lambda f.f A t$ and $\text{open } t \text{ as } \langle Z, y \rangle \text{ in } s = t V (\Lambda Z.\lambda y.s)$. Using parametricity we can prove the following commutation property and η -rule for existential types:

Lemma 5.8 *Assume the underlying bifibration satisfies the Beck-Chevalley condition, and that Eq is full.*

- (i) Let $\Gamma; \Delta \vdash t : \exists X.T$, let $\Gamma, Z; \Delta, u : T[Z/X] \vdash s : S$ and let $\Gamma; \Delta \vdash f : S \rightarrow S'$ for a closed type S' . Then $\llbracket f(\text{open } t \text{ as } \langle Z, u \rangle \text{ in } s) \rrbracket_o = \llbracket \text{open } t \text{ as } \langle Z, u \rangle \text{ in } f(s) \rrbracket_o$.
- (ii) If $\Delta; \Gamma \vdash t : \exists X.T$, then $\llbracket \text{open } t \text{ as } \langle Z, u \rangle \text{ in } \langle Z, u \rangle \rrbracket_o = \llbracket t \rrbracket_o$. \square

If $F : \mathcal{C} \rightarrow \mathcal{C}$ is an endofunctor, an F -coalgebra is a pair (A, k_A) with A an object of \mathcal{C} and $k_A : A \rightarrow FA$ a morphism. We call A the *carrier* of the F -coalgebra and k_A its *structure map*. A morphism $h : A \rightarrow B$ in \mathcal{C} is an F -coalgebra homomorphism $h : (A, k_A) \rightarrow (B, k_B)$ if $k_B \circ h = Fh \circ k_A$. An F -coalgebra (W, out) is *weakly final* if, for any F -coalgebra (A, k_A) , there exists a mediating F -coalgebra homomorphism $unfold[A, k_A] : (A, k_A) \rightarrow (W, out)$. It is a *final F -coalgebra* if $unfold[A, k_A]$ is unique.

Let $F = (F_o, F_r) : \text{Rel}(U) \rightarrow_{\text{Eq}} \text{Rel}(U)$ be a functor with a strength t . We show that the final F_o -coalgebra exists. Again, we first construct a weakly final coalgebra

by defining $W = (W_o, W_r) = [\exists X.(X \rightarrow \underline{F}(X)) \times X]$.

Lemma 5.9 W_o is the carrier of a weakly final F_o -coalgebra (W_o, out_o) with mediating morphism $unfold_o[A, k]$ and W_r is the carrier of a weakly final F_r -coalgebra (W_r, out_r) with mediating morphism $unfold_r[A, k]$. \square

We proceed similarly to Lemma 5.6. This time, we use the opfibrational part of the Graph Lemma to construct F_r -coalgebras.

Lemma 5.10 Assume the underlying bifibration satisfies the Beck-Chevalley condition, and that Eq is full.

- (i) For every F_o -coalgebra morphism $h : (A, k_A) \rightarrow (B, k_B)$ we have $unfold_o[B, k_B] \circ h = unfold_o[A, k_A]$.
- (ii) $unfold_o[W_o, out_o] = id_{W_o}$. \square

Putting things together, we have constructed a final coalgebra.

Theorem 5.11 If the underlying bifibration satisfies the Beck-Chevalley condition, and if Eq is full, then (W_o, out_o) is a final F_o -coalgebra. \square

5.4 Parametricity Implies Dinaturality

We show that our axiomatic foundations can be used to prove that dinaturality can be deduced from parametricity. This is well-known in other settings (see, e.g., [4, Section 5.1]), but we do it because i) it shows our foundation passes this test; and ii) it highlights again the use of bifibrations to give two definitions of the graph of a function both of which are used in the proof. First, the definition of dinaturality:

Definition 5.12 If $F, G : \mathcal{B}^{op} \times \mathcal{B} \rightarrow \mathcal{B}$ are mixed variant functors, then a dinatural transformation $t : F \rightarrow G$ is a collection of morphisms $t_X : FXX \rightarrow GXX$ indexed by objects X of \mathcal{B} such that, for every $g : X \rightarrow Y$ of \mathcal{B} , the following commutes:

$$\begin{array}{ccc}
 & F X X & \xrightarrow{t_X} G X X \\
 F(g, id_X) \nearrow & & \searrow G(id_X, g) \\
 F Y X & & G X Y \\
 F(id_Y, g) \searrow & & \nearrow G(g, id_Y) \\
 & F Y Y & \xrightarrow{t_Y} G Y Y
 \end{array}$$

We note that our proof applies to all mixed variant functors with equality preserving liftings, not just strong such functors.

Theorem 5.13 Let $(F_o, F_r), (G_o, G_r) : \text{Rel}(U)^{op} \times \text{Rel}(U) \rightarrow_{\text{Eq}} \text{Rel}(U)$. Further, let $t_A^0 : F_oAA \rightarrow G_oAA$ be a family indexed by objects A of \mathcal{B} , and $t_R^1 : F_rRR \rightarrow G_rRR$ be a family indexed by objects R of $\text{Rel}(\mathcal{E})$ such that if R is over (A, B) , then t_R^1 is over (t^0A, t^0B) . Then t^0 is a dinatural transformation from F_o to G_o . \square

Theorem 5.13 applies in particular to the interpretation of terms $t : \forall X.FXX \rightarrow GXX$ where F and G are given by type expressions with two free type variables, one occurring positively and one negatively. As is well known, dinaturality reduces to naturality when F and G are covariant.

6 Examples

The construction of examples remains delicate — for instance, there are no set-theoretic models with a classical meta-theory. We give five models: Examples 6.1, 6.3, 6.4 and 6.5 are to be regarded as being internal to the Calculus of Constructions with impredicative **Set** (with $\neg\neg$ -stable equality for Example 6.3), while Example 6.2 is internal to the category of ω -sets.

Before doing so, we take a moment to emphasise the generality of our framework. Considering different fibrations, we can derive parametric models with very different flavours. For example, changing the base category of the fibration corresponds to changing the ‘standard’ model in which we interpret types and terms. Changing the total category and the fibration (i.e., the functor itself) corresponds to changing the relevant notion of relational logic. We take advantage of the possibility of non-standard relations in Examples 6.2, 6.3 and Non-example 6.5.

Example 6.1 Reynolds’ set-theoretic model is an instance of our framework via the relations fibration on **Set**. The equality functor is full and faithful in this bifibration, and **Set** is well-pointed. Hence Theorems 5.7 and 5.13 ensure that initial algebras exist, and that all terms are interpreted as dinatural transformations.

Example 6.2 The PER model of Bainbridge et al. [2] is an instance of our framework, if bifibrations are understood as internal to the category of ω -sets, so that natural transformations are *uniformly* realised (see also Longo and Moggi [18] for a detailed construction of a model using a category of PERs internal to ω -sets). An object of the category $\text{PER}_{\mathbb{N}}$ is a symmetric, transitive relation $R \subseteq \mathbb{N} \times \mathbb{N}$. A morphism from R to S is a function $f : \mathbb{N}/R \rightarrow \mathbb{N}/S$ that is *tracked* by some partial recursive function $\phi_k : \mathbb{N} \rightarrow \mathbb{N}$, i.e., such that $f([n]_R) = [\phi_k(n)]_S$ for all $[n]_R \in \mathbb{N}/R$. The appropriate notion of predicate with respect to a PER R is that of a *saturated* subset, i.e., a subset $P \subseteq \mathbb{N}$ such that $P(x)$ and $R(x, x')$ implies $P(x')$. Saturated subsets form a bifibration over PERs with a full and faithful equality functor $\text{Eq}A = A$. The CCC structure of $\text{PER}_{\mathbb{N}}$ and **SatRel** is standard; a bijective pairing function $\langle -, - \rangle : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$ gives the product and recursion theory (the *s-m-n* Theorem) gives the exponential. The interesting case is that of forall types, which are interpreted as (cut-down, to ensure equality preservingness) intersections of PERs: $\llbracket \forall X.F \rrbracket_o \bar{R} = \{(n, k) \in \bigcap_{R' : \text{PER}_{\mathbb{N}}} \llbracket F \rrbracket_o(\bar{R}, R') \mid \text{for all } Q : \text{SatRel}(S, T). (n, n), (k, k) \in \llbracket F \rrbracket_r(\text{Eq}(\bar{R}), Q)\}$ and $\llbracket \forall X.F \rrbracket_r \bar{P} = \bigcap_{Q : \text{SatRel}(R, S)} \llbracket F \rrbracket_r(\bar{P}, Q)$. Since $\text{PER}_{\mathbb{N}}$ is also well-pointed, Theorems 5.7 and 5.13 again apply.

Example 6.3 The previous models are well-known, but our framework also suggests new ones. A relation $R \subseteq X \times Y$ can be understood classically as a function from $X \times Y$ to **Bool**. (Constructively, this only covers *decidable* relations.) Here, **Bool** can be replaced with any constructively completely distributive [8] non-trivial lattice \mathcal{V} of “truth values”, leading to “multi-valued parametricity”. For instance, the collection $\mathcal{D}(L)$ of all down-closed subsets of a complete lattice L is constructively completely distributive, and classically, we recover **Bool** as $\mathcal{D}(\mathbf{1})$. The category $\text{Fam}(\mathcal{V})$ has objects (A, p) , where A is a set and $p : A \rightarrow \mathcal{V}$ is thought of as a \mathcal{V} -valued predicate. The families fibration $\pi : \text{Fam}(\mathcal{V}) \rightarrow \text{Set}$ is a bifibration with $\Sigma_f(Q)(y) = \sup_{f x=y} Q(x)$, fibred terminal objects $(X, \lambda _ . \top)$, where \top is the greatest

element of \mathcal{V} , and comprehension given by $\{(A, p)\} = p^{-1}(\top)$. Since \mathcal{V} is complete, it is a Heyting algebra, so that $\pi : \mathbf{Fam}(\mathcal{V}) \rightarrow \mathbf{Set}$ is a fibred CCC. Also, π has simple products given by $\Pi_\pi(A \times B, p)(a) = \inf_{x \in B} p(a, b)$. By Lemma 4.4, $\mathbf{Rel}(\pi)$ is thus an equality preserving arrow fibration. Finally, the interpretation of forall types is given by $\llbracket \forall X.F \rrbracket_o \bar{A} = \{f : \prod_{S:\mathbf{Set}} \llbracket F \rrbracket_o(\bar{A}, S) \mid \inf_{P:X \times Y \rightarrow \mathcal{V}} \llbracket F \rrbracket_r(\mathbf{Eq}\bar{A}, P) = \top\}$ and $\llbracket \forall X.F \rrbracket_r \bar{P} = \inf_{Q:X \times Y \rightarrow \mathcal{V}} \llbracket F \rrbracket_r(P, Q)$. Distributivity is used to show that this functor is equality preserving. Fullness of \mathbf{Eq} is obvious by $\neg\neg$ -stable equality.

The extra conditions we impose in Section 5 really are necessary: the following are examples of \forall - and arrow-fibrations where Theorems 5.7 and 5.13 fail.

Non-example 6.4 Let G be a fixed (non-trivial) group, and consider the category of G -sets, i.e., the category with objects (X, \cdot_X) , where X is a set and $\cdot_X : G \times X \rightarrow X$ is a group action. Morphisms are functions between the carrier sets that respect the group action. Together with equivariant (i.e., group action respecting) relations, this forms a bifibration that is a model of System F in the sense of Theorem 4.6. However, the category is not well-pointed, and we can see that this makes Theorem 5.7 fail in our setting: the interpretation of $\forall X.X \rightarrow X$ is not the singleton G -set $\mathbf{1}$ as expected, but instead contains all the elements of the group G . We conjecture that this non-example also extends to a constructive treatment of the category of nominal sets [23].

Non-example 6.5 The identity fibration $\mathbf{Id} : \mathbf{Set} \rightarrow \mathbf{Set}$ models *ad hoc* polymorphism: it is a \forall - and arrow-fibration, but the equality functor $\mathbf{Eq} X = X \times X$ is not full. This explains why Theorem 5.13 fails: $\llbracket \forall X.T \rrbracket_o$ includes *ad hoc* polymorphic functions, so that e.g. $\llbracket \forall X.X \rightarrow X \rrbracket_o$ contains non-natural transformations such as η , where $\eta_{\mathbf{Bool}}(x) = \neg x$ and $\eta_X(x) = x$ for $X \neq \mathbf{Bool}$.

7 Conclusions and future work

Our interpretation of types and terms as fibred functors and fibred natural transformations shows that parametricity entails replacing the usual categorical semantics involving categories, functors, and transformations with one based on fibrations, fibred functors, and fibred transformations. The results in Section 5 show that our new approach based on *bifibrations* hits the sweet spot of a light structure that still suffices to prove key results. Work is ongoing in using the bifibrational framework to develop new notions such as proof-relevant parametricity, and higher order parametricity with interesting links to cubical sets that also appear in the semantics of Homotopy Type Theory [3].

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